

Corrections

Corrigendum to "Line Digraph Iterations and Connectivity Analysis of de Bruijn and Kautz Graphs"

Ding-Zhu Du, D. Frank Hsu, and Yuh-Dauh Lyuu

An error appears in the proof of Lemma 3.2 in our paper "Line Digraph Iterations and Connectivity Analysis of de Bruijn and Kautz Graphs" [Ding-Zhu Du, Yuh-Dauh Lyuu, and D. Frank Hsu, *IEEE Trans. Computers*, vol. 42, no. 5, pp. 612-616, May 1993]. To fix it, we need to modify Definition 3.1 and the Proof of Lemma 3.2 on p. 614. They should read as follows.

DEFINITION 3.1. Given positive integers k and l , let $\mathcal{L}(k, l)$ be the class of graphs G so that for any $m \leq k$ distinct nodes $y_1, \dots, y_m \in V(G)$, a node $x \in V(G)$, any $2q$ ($1 \leq q \leq k + 1$) edges $e_0, e_1, \dots, e_{2q-1}$ where $e'_0 = e'_2 = \dots = e'_{2q-1} = x''$ and

$$\{e''_1, e''_3, \dots, e''_{2q-1}\} \subseteq \{y'_1, y'_2, \dots, y'_m\},$$

and m positive numbers r_1, \dots, r_m such that $\sum_i r_i = k - q + 1$, there exist $k - q + 1$ node-disjoint paths each of positive length at most l from x to y_1, \dots, y_m , where r_i of them go to y_i for $1 \leq i \leq m$, and none involves e_j for $0 \leq j \leq 2q - 1$. Note: x may be one of the y_i . (The introduction calls this notion spread*.)

LEMMA 3.2. $L^h(G) \in \mathcal{L}(k, l + h)$ if $G \in \mathcal{L}(k, l)$, for $h \geq 0$.

PROOF. The proof is by induction on i , the depth of line graph iterations. This lemma surely holds for $i = 0$. Assume it is true for i up to $h - 1$ and consider $i = h > 0$.

Consider a node x and m distinct nodes y_1, \dots, y_m in $L^i(G)$. These nodes are edges in $L^{i-1}(G)$. Let r_1, \dots, r_m be positive integers such that $\sum_j r_j = k - 1 + 1$ ($1 \leq q \leq k + 1$). Assume $x'' = y'_1 = \dots = y'_s$ in $L^{i-1}(G)$ for some $0 \leq s \leq m$, with $s = 0$ meaning no (x, y_j) is an edge in $L^i(G)$. Hence, in $L^{i-1}(G)$, the edge x is incident to edges y_j for $1 \leq j \leq s$.

There are three cases to consider:

- 1) $s = 0$,
- 2) $s > 0$ and $\{(x, y_j) \mid 1 \leq j \leq s\} \subseteq \{e_0, e_1, \dots, e_{2q-1}\}$, and
- 3) $s > 0$ and there exists $1 \leq j \leq s$ such that $(x, y_j) \notin \{e_0, e_1, \dots, e_{2q-1}\}$.

First, suppose 1) occurs, i.e., $s = 0$. By the induction hypothe-

sis, in $L^{i-1}(G)$, there are $k - q + 1$ node-disjoint paths from x'' to y'_1, \dots, y'_m where r_j of them end at y'_j and each has positive length at most $l + i - 1$ (here we treat y'_a and y'_b with $a \neq b$ as distinct even if $y'_a = y'_b$) and none involves $2q$ edges $e''_0, e''_1, e''_2, \dots, e''_{2q-1}$, in $L^{i-1}(G)$. Typical paths look like

$$\overbrace{x'' \rightarrow \dots \rightarrow y'_j}^{\alpha}.$$

The paths extended to

$$x' \rightarrow \overbrace{x'' \rightarrow \dots \rightarrow y'_j}^{\alpha} \rightarrow y''_j.$$

they induce, in $L^i(G)$, the desired node-disjoint paths from x to the y'_j 's.

Suppose 2) occurs. By the same argument as the above, we can prove the existence of the desired node-disjoint paths from x to y_j 's.

Now, suppose 3) occurs. Without loss of generality, assume $(x, y_1) \notin \{e_0, e_1, \dots, e_{2q-1}\}$. If $q = k + 1$, then the desired $k - q + 1$ paths exist trivially. Thus, we may also assume $q < k + 1$. Define $y_{m+1} = y_1$ with $r_{m+1} = r_1 - 1$ (if $r_{m+1} = 0$, just drop node y_{m+1} from the induction step). Consider the graph obtained from $L^{i-1}(G)$ by deleting edges x, y_1 , and $2q$ edges $e''_0, e''_1, e''_2, \dots, e''_{2q-1}$ and considering $y'_{s+1}, \dots, y'_{m+1}$. By the induction hypothesis, there are $\sum_{j=2}^{m+1} r_j = k - q$ node-disjoint paths in $L^{i-1}(G)$ from x'' to y'_2, \dots, y'_{m+1} where r_j of them end at y'_j (here we treat y'_a and y'_b with $a \neq b$ as distinct even if $y'_a = y'_b$), each has positive length at most $l + i - 1$, and none involves x, y_1 , and $2q$ edges $e''_0, e''_1, e''_2, \dots, e''_{2q-1}$. Typical paths look like

$$\overbrace{x'' \rightarrow \dots \rightarrow y'_j}^{\alpha} \text{ for } 2 \leq j \leq m + 1.$$

The paths extend to

$$x' \rightarrow \overbrace{x'' \rightarrow \dots \rightarrow y'_j}^{\alpha} \rightarrow y''_j, \text{ for } 2 \leq j \leq m + 1,$$

they induce, in $L^i(G)$, $k - q$ node-disjoint paths from x to y_j 's, each of positive length at most $l + i$ and none involving the edge (x, y_1) and $2q$ edges $e_0, e_1, \dots, e_{2q-1}$. Add (x, y_1) to get the $(k - q + 1)$ th path in $L^i(G)$ and the proof is completed. \square

With the above modification, the remainder in the paper is correct except that in the introduction, the phrase "two consecutive edges" at line 28 of p. 613 should be replaced by "some edges."

ACKNOWLEDGMENT

The authors are indebted to C. Padró, P. Morillo, and M.A. Fiol for pointing out this error.

- D.-Z. Du is with the Department of Computer Science, University of Minnesota, Minneapolis, MN 55455. E-mail: dzd@cs.umn.edu.
- D.F. Hsu is with the Department of Computer and Information Science, Fordham University, Bronx, NY 10458.
- Y.-D. Lyuu is with the Department of Computer Science, University of Taiwan, Taipei, Taiwan..

Manuscript received Jan. 18, 1996.

For information on obtaining reprints of this article, please send e-mail to: transcom@computer.org, and reference IEEECS Log Number C96105.